

A bipartite graph  $G = (V^+, V^-; A)$  with  $V^+ \cup V^- \neq \emptyset$  is said to be *DM-irreducible* if it cannot be decomposed into more than one nonempty component in the DM-decomposition. Otherwise, it is called *DM-reducible*. A graph  $G$  with  $V^+ = \emptyset$  or  $V^- = \emptyset$  is DM-irreducible, as the whole graph is a (vertical or horizontal) tail. Note that  $G$  with  $A = \emptyset$ ,  $V^+ \neq \emptyset$  and  $V^- \neq \emptyset$  is DM-reducible, as it can be decomposed into two nonempty components, the horizontal tail  $G_0 = (V^+, \emptyset; \emptyset)$  and the vertical tail  $G_\infty = (\emptyset, V^-; \emptyset)$ .

The following theorem is a reformulation of the result due to Marcus-Minc [186] and Brualdi [22] (see also Brualdi-Ryser [24, Theorem 4.2.2]).

**Theorem 2.2.24.** For a bipartite graph  $G = (V^+, V^-; A)$  with  $|V^+| = |V^-|$  the following three conditions are equivalent:

- (i)  $G$  is DM-irreducible.
- (ii)  $C(G) = \{(V^+, \emptyset), (\emptyset, V^-)\}$ .
- (iii)  $\nu(G \setminus \{u, v\}) = \nu(G) - 1$  for  $\forall u \in V^+, \forall v \in V^-$ .

*Proof.* The equivalence between (i) and (ii) follows from Theorem 2.2.22(1). For (ii)  $\Rightarrow$  (iii), first note (ii) implies  $\nu(G) = |V^+|$ , and hence  $\nu(G \setminus \{u, v\}) \leq |V^+| - 1 = \nu(G) - 1$ . Take a minimum cover  $(W^+, W^-)$  of  $G \setminus \{u, v\}$ . Since  $(W^+ \cup \{u\}, W^- \cup \{v\})$  is a cover of  $G$  but not a minimum cover, we see  $\nu(G) + 1 \leq |W^+ \cup \{u\}| + |W^- \cup \{v\}| = \nu(G \setminus \{u, v\}) + 2$ .

For (iii)  $\Rightarrow$  (ii) suppose that (ii) fails. We divide into two cases: (a)  $\nu(G) = |V^+|$  and (b)  $\nu(G) < |V^+|$ . In case (a), there exists  $(U^+, U^-) \in C(G)$  such that  $U^+ \neq \emptyset$  and  $U^- \neq \emptyset$ . For  $u \in U^+$  and  $v \in U^-$ ,  $(U^+ \setminus \{u\}, U^- \setminus \{v\})$  is a cover of  $G \setminus \{u, v\}$ . Hence  $\nu(G \setminus \{u, v\}) \leq |U^+ \setminus \{u\}| + |U^- \setminus \{v\}| = \nu(G) - 2$ . In case (b), both horizontal and vertical tails are nonempty. For  $u \in V_0^+$  and  $v \in V_\infty^-$  we have  $\nu(G \setminus \{u, v\}) = \nu(G)$  by Corollary 2.2.23(2).  $\blacksquare$

The concept of DM-decomposition may be extended to matrices by means of the DM-decomposition of associated bipartite graphs. Recall from §2.2.1 that for a matrix  $A = (A_{ij})$  the associated bipartite graph is defined by  $G = (V^+, V^-; A)$  with  $V^+ = \text{Col}(A)$ ,  $V^- = \text{Row}(A)$  and  $\bar{A} = \{(j, i) \mid A_{ij} \neq 0\}$ . A DM-component  $G_k = (V_k^+, V_k^-; A_k)$  of  $G$  corresponds to the submatrix  $A[V_k^-, V_k^+]$ , which will be referred to as a *DM-component* of  $A$ .

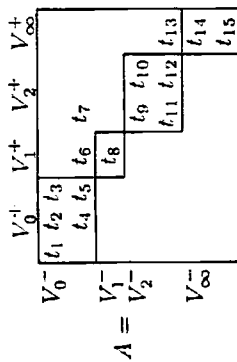
The following relation is obvious from the definitions, but provides the DM-decomposition with a linear algebraic significance.

**Proposition 2.2.25.** term-rank  $A = \nu(G)$ .  $\square$

The *DM-decomposition* of  $A$  gives the finest block-triangularization of a matrix by means of a transformation  $P_r A P_c$  using two permutation matrices  $P_r$  and  $P_c$ , where it is imposed that each diagonal block in a block-triangularization has full term-rank. In fact, Theorem 2.2.22(1) combined with Proposition 2.2.25 above guarantees this term-rank condition for the diagonal blocks produced by the DM-decomposition. Furthermore, term-rank

coincides with rank for a generic matrix, in which all nonzero entries are independent parameters (cf. Proposition 2.1.12). Hence, for a generic matrix, the DM-decomposition gives the finest proper block-triangularization in the sense of §2.1.4.

For instance, the matrix version of the DM-decomposition of the graph in Fig. 2.7 is given by



Note that term-rank  $A[V_k^-, V_k^+] = \min(|V_k^+|, |V_k^-|)$  for  $k = 0, 1, 2, \infty$ . The matrix  $P_r A P_c$  of (2.16) in Example 2.1.15 gives an instance of the DM-decomposition of a term-nonsingular matrix.

Though term-rank is a natural combinatorial counterpart, it is not the same as rank, which is undoubtedly more important in applications. A numerical (nongeneric) matrix may or may not have the same rank as term-rank, and accordingly, the DM-decomposition may or may not be a proper block-triangular form. The present argument shows the following.

**Proposition 2.2.26.** For a matrix  $A$  the following three conditions (i)-(iii) are equivalent.

- (i) rank  $A = \text{term-rank } A$ .
- (ii) The DM-decomposition is a proper block-triangularization, i.e., the DM-components  $A[V_k^-, V_k^+]$  ( $k = 0, 1, \dots, b, \infty$ ) satisfy

$$\text{rank } A[V_k^-, V_k^+] = \min(|V_k^+|, |V_k^-|) \quad (k = 0, 1, \dots, b, \infty).$$

- (iii) There exist  $I \subseteq \text{Row}(A)$  and  $J \subseteq \text{Col}(A)$  such that rank  $A[I, J] = 0$ , rank  $A[\text{Row}(A) \setminus I, J] = |\text{Row}(A) \setminus I|$ , and rank  $A[I, \text{Col}(A) \setminus J] = |\text{Col}(A) \setminus J|$ .

*Proof.* The equivalence of (i) and (ii) is immediate from Theorem 2.2.22. For (iii) take  $I = \text{Row}(A) \setminus V_0^-$  and  $J = \text{Col}(A) \setminus V_0^+$ .  $\blacksquare$

The concept of DM-irreducibility can be naturally defined for matrices, and it coincides with the well-studied concept of full indecomposability (cf. Brualdi-Ryser [24] and Schneider [288]). To see this, first recall that a matrix  $A$  is said to be *fully indecomposable* if it does not contain a zero submatrix  $A[I, J] = O$  with  $I \neq \emptyset$ ,  $J \neq \emptyset$  and  $|I| + |J| = \max(|\text{Row}(A)|, |\text{Col}(A)|)$ . Since  $A[I, J] = O$  if and only if  $(V^- \setminus I, V^+ \setminus J)$  is a cover of the associated graph  $G = (V^+, V^-; A)$ , matrix  $A$  is fully indecomposable if and only if  $G$  has no cover  $(U^+, U^-)$  such that  $U^+ \neq V^+$ ,  $U^- \neq V^-$  and

$|U^+| + |U^-| = \min(|V^+|, |V^-|)$ . The latter condition is equivalent to the DM-irreducibility (cf. Theorem 2.2.24). Henceforth we use DM-irreducibility as a synonym of full indecomposability.

The DM-irreducibility for square generic matrices admits two further characterizations in addition to those given in Theorem 2.2.24. The first says that the DM-irreducibility for a generic matrix is equivalent to the inverse matrix having a completely dense nonzero pattern.

**Theorem 2.2.27.** *A square generic matrix  $A$  is DM-irreducible if and only if  $A$  is nonsingular and  $(A^{-1})_{ji} \neq 0$  for all  $(j, i)$ .*

*Proof.* Since  $(A^{-1})_{ji} = \det A[R \setminus \{i\}, C \setminus \{j\}] / \det A$ , where  $R = \text{Row}(A)$  and  $C = \text{Col}(A)$ , the claim here reduces to the equivalence of (i) and (iii) in Theorem 2.2.24. ■

The determinant of a generic matrix  $A$  can be regarded as a polynomial in the nonzero entries. Specifically, let  $\mathcal{T}$  denote the set of nonzero entries of  $A$ , which is algebraically independent over a ground field  $K$ . Then  $\det A \in K[\mathcal{T}]$ , where  $K[\mathcal{T}]$  means the ring of polynomials in  $\mathcal{T}$  over  $K$ .

The following theorem gives an algebraic characterization of the DM-irreducibility in terms of the irreducibility of the determinant as a multivariate polynomial. This is proven in Ryser [285] and credited essentially to Frobenius [78] in Ryser [286].

**Theorem 2.2.28.** *A square generic matrix  $A$  is DM-irreducible if and only if  $\det A$  is an irreducible (nonzero) polynomial in  $K[\mathcal{T}]$ , where  $\mathcal{T}$  denotes the set of nonzero entries of  $A$ .*

*Proof.* The “if” part is obvious, since, for a DM-reducible  $A$  with no tails,  $\det A$  is equal to the product of the determinants of the diagonal blocks of the DM-decomposition of  $A$  (and  $\det A = 0$  if a nonempty tail exists). For the “only if” part assume that  $\det A$  is factored as  $\det A = f_1 \cdot f_2$  with  $f_1, f_2 \in K[\mathcal{T}] \setminus K$ . For  $k = 1, 2$ , let  $\mathcal{T}_k$  denote the set of the variables of  $\mathcal{T}$  that appear in  $f_k$ . Put

$$R_k = \{i \in R \mid A_{ij} \in \mathcal{T}_k\}, \quad C_k = \{j \in C \mid A_{ij} \in \mathcal{T}_k\} \quad (k = 1, 2),$$

where  $R = \text{Row}(A)$  and  $C = \text{Col}(A)$ . Then  $R_1 \cap R_2 = \emptyset$ ,  $R_1 \cup R_2 = R$ ,  $C_1 \cap C_2 = \emptyset$ ,  $C_1 \cup C_2 = C$  for  $k = 1, 2$ , since for each pair of terms in  $f_1$  and  $f_2$  their product remains in  $f_1 \cdot f_2 = \det A$  as a nonvanishing term, which in turn corresponds to a perfect matching in the associated bipartite graph. We may assume  $|R_1| \geq |C_1| \geq 1$  without loss of generality. If  $A[R_1, C_2] = O$ ,  $A$  is DM-reducible. If  $A_{ij} \neq 0$  for some  $i \in R_1$  and  $j \in C_2$ , the variable  $t = A_{ij}$  cannot appear in  $\det A$ , since otherwise  $t$  must be contained in  $f_k$  for  $k = 1$  or  $2$ , which implies  $i \in R_k$  and  $j \in C_k$ , a contradiction to  $R_1 \cap R_2 = \emptyset$  and  $C_1 \cap C_2 = \emptyset$ . The disappearance of  $t$  in  $\det A$  implies  $\det A[R \setminus \{i\}, C \setminus \{j\}] = 0$ , or equivalently,  $\nu(G \setminus \{i, j\}) < \nu(G) - 1$  in terms of the associated bipartite graph  $G$ . This shows the DM-irreducibility by Theorem 2.2.24. ■

**Remark 2.2.29.** We have derived the DM-decomposition in a systematic manner on the basis of the Jordan–Hölder-type theorem for submodular functions, though alternative quicker derivations would have been possible. Our systematic derivation here enables us to generalize the DM-decomposition to a more sophisticated decomposition in §4.4, called the CCF (combinatorial canonical form) of layered mixed matrices. The DM-decomposition serves as one of the main tools for the graph-theoretic methods for systems analysis (see Duff–Erisman–Reid [59], Murota [204, Chaps. 2 and 3]), whereas the CCF is for the matroid-theoretic methods to be developed in Chap. 4 and Chap. 6. Applications of the DM-decomposition can be found in Ashcraft–Liu [8], Erisman–Grimes–Lewis–Poole–Simon [73], Hellerman–Rarick [109, 110], O’Neil–Szyld [255], and Pothen–Fan [273]. □

## 20 Algorithms and Combinatorics

# Matrices and Matroids for Systems Analysis

Kazuo Murota